
Anti-isomorphisms, character modules, and self-dual codes over non-commutative rings

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Abstract: This paper is dedicated to Vera Pless. It is an elaboration on ideas of Nebe, Rains, and Sloane: by assuming the existence of an anti-isomorphism on a finite ring and by assuming a module alphabet has a well-behaved duality, one is able to study self-dual codes defined over alphabets that are modules over a non-commutative ring. Various examples are discussed.

Keywords: anti-isomorphisms, bi-additive forms, character modules, dual codes, MacWilliams identities, skew-polynomial rings, group rings, Steenrod algebra.

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Biographical notes: Jay A. Wood is a professor of mathematics at Western Michigan University.

1 Tribute to Vera Pless

I first met Vera Pless on 12 October 1987 at the University of Wisconsin. She was visiting Madison at the time, and I was in town to give a talk. We had shared some correspondence prior to this, because I had discovered that a problem I had been pursuing, which originated in differential geometry and algebraic topology, turned out to be equivalent to classifying the doubly-even self-dual binary codes up to code equivalence. Vera very graciously provided me with copies of some of her papers on the classification of these codes, and I went on to study the relationships between self-dual codes and certain structures in algebraic topology.

When I re-located to the midwest in 1990, Vera very kindly invited me to speak in her UIC seminar on several occasions. On one such occasion, 28 April 1992, Vera suggested that I re-examine the work of MacWilliams on code equivalence. Inspired

by Vera's interest, and by subsequent collaboration with Thann Ward, I spent much of the next seventeen years developing an understanding of code equivalence and the MacWilliams identities for linear codes defined over finite rings and finite modules, with a special emphasis on the role of Frobenius rings. I doubt that I would have worked on these topics if Vera had not suggested them. So, I can truthfully say: "Vera, I owe it all to you. Thank you!"

I find it fitting that this paper allows me to return to the topic of my first correspondence with Vera: self-dual codes, with some examples coming from algebraic topology.

2 Introduction

My goal in this paper is to clarify the assumptions that lead a good theory of self-dual codes over non-commutative rings. The paper was inspired by, and is an elaboration of portions of, the book by Nebe, Rains, and Sloane (NRS06). Virtually all of the content of Sections 3 and 4 can be found, explicitly or implicitly, in (NRS06). The material in those sections has been organized in a manner similar to the latter portions of (Woo09), in which a detailed examination of duality and the MacWilliams identities took place.

In (Woo09), the dual of a left linear code is a right linear code, and vice versa. While the theory works well, the change of sides makes it difficult for a code to be self-dual. This problem is addressed in (NRS06), and in Section 3 the approach of (NRS06) is condensed into three properties (Definition 3.3): that the ring R admit an anti-isomorphism ε ; that the alphabet A , a finite left R -module, admit an anti-isomorphism ψ to its character module \hat{A} ; and that ψ have a close relation to its own dual map. These properties are equivalent to the existence of a bi-additive form $\beta : A \times A \rightarrow \mathbb{Q}/\mathbb{Z}$, with certain additional properties. The form β plays a major role in (NRS06), and the same is true here.

In Section 4 it is shown (Theorem 4.1) that the three properties of Definition 3.3 imply that left dual codes of left linear codes exist and have all the good properties one would expect of a dual code. In subsequent sections, the paper addresses each of the three properties in turn. In some cases, general theorems can be proved concerning when the property holds. In all cases, several examples (including a group ring and a finite subalgebra of the Steenrod algebra) are explored in detail.

3 Anti-isomorphisms and character modules

Let R be a finite ring with 1. We allow R to be non-commutative. Let A be a finite left R -module, which will serve as the alphabet for linear codes over R . A left R -linear code over A of length n is a left R -submodule $C \subset A^n$. (There is a parallel theory for right linear codes.) An important special case is when the alphabet A is the ring R itself, viewed as a left R -module.

In most treatments of the MacWilliams identities, the dual code C^\perp would be a right R -module. Unless the ring R is commutative, this change of sides would make it impossible for a code to be self-dual, i.e., satisfy $C = C^\perp$. Nebe, Rains, and Sloane (NRS06) address this problem by assuming some additional structure

on R and A so that one can view the dual code C^\perp as a left R -module. To that end, we follow (NRS06) and introduce several definitions.

An *anti-isomorphism* $\varepsilon : R \rightarrow R$ of a ring R is an isomorphism of abelian groups with the property that $\varepsilon(rs) = \varepsilon(s)\varepsilon(r)$ for all $r, s \in R$. An anti-isomorphism ε defines an isomorphism $R \cong R^{op}$ between the ring R and its opposite ring R^{op} . If ε is an anti-isomorphism of R , then so is its inverse ε^{-1} . An *involution* is an anti-isomorphism $\varepsilon : R \rightarrow R$ such that ε^2 is the identity; i.e., $\varepsilon^{-1} = \varepsilon$.

Let ${}_R\mathcal{F}$ (resp., \mathcal{F}_R) denote the category of finitely-generated left (resp., right) R -modules and R -module homomorphisms. Then an anti-isomorphism $\varepsilon : R \rightarrow R$ induces covariant functors $\varepsilon : {}_R\mathcal{F} \rightleftharpoons \mathcal{F}_R$ as follows. If M is a left R -module, define $\varepsilon(M)$ to be the same abelian group as M with right scalar multiplication defined by $mr = \varepsilon(r)m$, for $m \in M, r \in R$, where $\varepsilon(r)m$ uses the left scalar multiplication of M . Similarly, if N is a right R -module, then $\varepsilon(N)$ has left scalar multiplication defined by $rn = n\varepsilon(r)$, for $n \in N, r \in R$. One verifies that a homomorphism $f : M_1 \rightarrow M_2$ of left R -modules is also a homomorphism $\varepsilon(M_1) \rightarrow \varepsilon(M_2)$ of right R -modules.

Character modules will be important in our discussion, so we provide a short summary of them next. The *character functor* $\widehat{} : {}_R\mathcal{F} \rightleftharpoons \mathcal{F}_R$ is a contravariant functor that associates to every finite left (resp., right) R -module M its character module $\widehat{M} = \text{Hom}_{\mathbb{Z}}(M, \mathbb{Q}/\mathbb{Z})$, which is a finite right (resp., left) R -module. (In this paper, the additive form of characters will be used. By composing with the exponential map: $\mathbb{Q}/\mathbb{Z} \rightarrow \mathbb{C}^\times, x \mapsto \exp(2\pi ix), x \in \mathbb{Q}/\mathbb{Z}$, one recovers the multiplicative form of characters. The modules involved are isomorphic.) When M is a left R -module, the right module structure of \widehat{M} is given by $(\varpi r)(m) = \varpi(rm)$, for $\varpi \in \widehat{M}, r \in R$, and $m \in M$.

Lemma 3.1. *Given an anti-isomorphism ε on a finite ring R , the functors ε and $\widehat{}$ commute. That is, for any finite R -module M ,*

$$\widehat{\varepsilon(M)} = \varepsilon(\widehat{M}).$$

Proof. When M is a left R -module, both $\varepsilon(M)$ and \widehat{M} are right R -modules, and $\widehat{\varepsilon(M)}$ and $\varepsilon(\widehat{M})$ are both left R -modules. For a character $\varpi : M \rightarrow \mathbb{Q}/\mathbb{Z}$, left multiplication by $r \in R$ in $\widehat{\varepsilon(M)}$ means, for $m \in M$, $(r\varpi)(m) = \varpi(mr)$, where the right multiplication is that of $\varepsilon(M)$. Left multiplication in $\varepsilon(\widehat{M})$ means $(r\varpi)(m) = (\varpi\varepsilon(r))(m)$. The two left multiplications agree because $\varpi(mr) = \varpi(\varepsilon(r)m) = (\varpi\varepsilon(r))(m)$, using the right module structures of $\varepsilon(M)$ and \widehat{M} . \square

Suppose the finite ring R admits an anti-isomorphism ε . Even though the functors ε and $\widehat{}$ commute, the functors cannot be the same. Indeed, the functor $\varepsilon : {}_R\mathcal{F} \rightleftharpoons \mathcal{F}_R$ is covariant, while the character functor $\widehat{} : {}_R\mathcal{F} \rightleftharpoons \mathcal{F}_R$ is contravariant. However, modules where the functors agree will be important.

To that end, suppose M is a finite left R -module such that $\psi : \varepsilon(M) \rightarrow \widehat{M}$ is an isomorphism of right R -modules. For every $y \in M$, $\psi(y)$ is a character on M . We denote the value of this character on a point $x \in M$ by $\psi(y)(x)$. Then ψ being a homomorphism means

$$\psi(\varepsilon(r)y)(x) = \psi(yr)(x) = (\psi(y)r)(x) = \psi(y)(rx), \quad (3.1)$$

for $r \in R$ and $x, y \in M$.

By applying the character functor to the isomorphism $\psi : \varepsilon(M) \rightarrow \widehat{M}$ and using that the double character module of M is naturally isomorphic to M itself, we obtain $\widehat{\psi} : M \rightarrow \widehat{\varepsilon(\widehat{M})} = \varepsilon(\widehat{M})$. Applying ε^{-1} , we have an isomorphism $\widehat{\psi} : \varepsilon^{-1}(M) \rightarrow \widehat{M}$. From the definition of $\widehat{\psi}$ we have the relation

$$\widehat{\psi}(x)(y) = \psi(y)(x), \quad x, y \in M. \quad (3.2)$$

Proposition 3.2. *Suppose a finite ring R admits an anti-isomorphism ε and that a finite left R -module M admits an isomorphism $\psi : \varepsilon(M) \rightarrow \widehat{M}$. Then $\varepsilon(M) \cong \varepsilon^{-1}(M)$; i.e., $\varepsilon^2(M) \cong M$.*

Proof. The composition $\widehat{\psi}^{-1}\psi : \varepsilon(M) \rightarrow \widehat{M} \rightarrow \varepsilon^{-1}(M)$ is an isomorphism. \square

Definition 3.3. The following is a list of properties that a ring R may possess.

P1: The ring R admits an anti-isomorphism ε .

P2: Given an anti-isomorphism ε on R , there exists a finite left R -module A and an isomorphism $\psi : \varepsilon(A) \rightarrow \widehat{A}$.

P3: In P2, the isomorphism ψ satisfies $\widehat{\psi} = \psi e$, for some unit $e \in R$.

The condition in P3 means that there exists a unit $e \in R$ such that $\widehat{\psi}(x) = \psi(x)e \in \widehat{A}$, for all $x \in A$, where $\psi(x)e$ uses the right module structure of \widehat{A} . This leads to the following relations:

$$\psi(x)(ey) = (\psi(x)e)(y) = \widehat{\psi}(x)(y) = \psi(y)(x), \quad x, y \in A. \quad (3.3)$$

For the rest of this section we assume that a finite ring R and a finite left R -module A satisfy P1–P3, with anti-isomorphism ε and right module isomorphism $\psi : \varepsilon(A) \rightarrow \widehat{A}$. In this case, observe that $\psi : A \rightarrow \varepsilon^{-1}(\widehat{A})$ is a left module isomorphism.

We will now associate to $\psi : \varepsilon(A) \rightarrow \widehat{A}$ a bi-additive form, in the spirit of (NRS06). First define $\beta : A \times A \rightarrow \mathbb{Q}/\mathbb{Z}$ by $\beta(a, b) = \psi(b)(a)$, for $a, b \in A$. Then extend β to $\beta : A^n \times A^n \rightarrow \mathbb{Q}/\mathbb{Z}$ by

$$\beta(x, y) = \sum_{i=1}^n \beta(x_i, y_i) = \sum_{i=1}^n \psi(y_i)(x_i), \quad (3.4)$$

for $x = (x_1, \dots, x_n), y = (y_1, \dots, y_n) \in A^n$. Remember that the notation $\psi(y_i)(x_i)$ means to evaluate the character $\psi(y_i)$ of A on the element $x_i \in A$. The result is an element of \mathbb{Q}/\mathbb{Z} . One then sums these elements of \mathbb{Q}/\mathbb{Z} .

Theorem 3.4. *Assume properties P1–P3. The form β of (3.4) satisfies:*

1. *The form β is bi-additive; i.e., $\beta(x+z, y) = \beta(x, y) + \beta(z, y)$ and $\beta(x, y+z) = \beta(x, y) + \beta(x, z)$, for all $x, y, z \in A^n$.*
2. *The form β is non-degenerate. That is, if $\beta(A^n, y) = 0$, then $y = 0$; and if $\beta(x, A^n) = 0$, then $x = 0$.*
3. *The form satisfies $\beta(rx, y) = \beta(x, \varepsilon(r)y)$, for all $r \in R, x, y \in A^n$.*

4. There exists a unit $e \in R$ so that $\beta(x, y) = \beta(ey, x)$, for all $x, y \in A^n$.

Conversely, assume property P1 and that there exists a form $\beta : A \times A \rightarrow \mathbb{Q}/\mathbb{Z}$ satisfying the properties above. If one defines $\psi : A \rightarrow \widehat{A}$ by $\psi(b)(a) = \beta(a, b)$, for $a, b \in A$, then ψ satisfies P2–P3.

Proof. The first part of the bi-additivity of β follows from the definition of character; the second from ψ being an additive homomorphism. The non-degeneracy of β follows from a similar non-degeneracy property of characters. The scalar multiplication property follows from ψ being a homomorphism of right R -modules and the definitions of the right module structures on \widehat{A} and $\varepsilon(A)$; compare with (3.1). The symmetry property follows from P3; compare with (3.3).

The converse is an exercise for the reader. \square

4 Dual codes and the MacWilliams identities

We first recall several well-known definitions. We assume that R is a finite ring with 1 and that A is a finite left R -module. An *additive code* of length n is an additive subgroup $C \subset A^n$. Recall that a left *linear code* is a left R -submodule of A^n . Every linear code is an additive code, but not conversely. The *Hamming weight* wt on A is a function $\text{wt} : A \rightarrow \mathbb{Q}$ defined by $\text{wt}(0) = 0$ and $\text{wt}(a) = 1$ for $a \neq 0$. The Hamming weight extends to a function $\text{wt} : A^n \rightarrow \mathbb{Q}$ by $\text{wt}(a_1, \dots, a_n) = \sum_{i=1}^n \text{wt}(a_i)$. The *Hamming weight enumerator* of an additive code $C \subset A^n$ is the polynomial $W_C(X, Y) \in \mathbb{C}[X, Y]$ defined by

$$W_C(X, Y) = \sum_{c \in C} X^{n-\text{wt}(c)} Y^{\text{wt}(c)}.$$

Assume that a finite ring R and a finite left R -module A satisfy P1–P3, with anti-isomorphism ε and right module isomorphism $\psi : \varepsilon(A) \rightarrow \widehat{A}$. The isomorphism ψ extends to a right module isomorphism $\psi : \varepsilon(A^n) \rightarrow \widehat{A}^n$.

Given an additive code $C \subset A^n$, the character-theoretic *annihilator* of C is

$$(\widehat{A}^n : C) = \{\varpi \in \widehat{A}^n : \varpi(C) = 0\}.$$

Note that $(\widehat{A}^n : C)$ is an additive subgroup of \widehat{A}^n and that $|C| |(\widehat{A}^n : C)| = |A^n|$; see (Woo99, (A.2)) or (Woo09, Section 2.2). Define the *dual code* C^\perp by

$$C^\perp = \psi^{-1}(\widehat{A}^n : C).$$

Note that the dual code C^\perp is an additive code in A^n . We say that C is *self-orthogonal* if $C \subset C^\perp$ and *self-dual* if $C = C^\perp$.

Theorem 4.1. *Assume that a finite ring R and a finite left R -module A satisfy P1–P3, with anti-isomorphism ε and right module isomorphism $\psi : \varepsilon(A) \rightarrow \widehat{A}$. Let β be the form associated to ψ via (3.4). Then:*

1. For any additive code $C \subset A^n$, $C^\perp = \{y \in A^n : \beta(C, y) = 0\}$.
2. For any additive code $C \subset A^n$, $|C| |C^\perp| = |A^n|$.



3. For any additive code $C \subset A^n$, the MacWilliams identities are satisfied:

$$W_{C^\perp}(X, Y) = \frac{1}{|C|} W_C(X + (|A| - 1)Y, X - Y).$$

4. If $C \subset A^n$ is a left linear code, then so is C^\perp .

5. For any left linear code $C \subset A^n$, $C = (C^\perp)^\perp$.

Proof. By the definition of C^\perp ,

$$C^\perp = \psi^{-1}(\widehat{A}^n : C) = \{y \in A^n : \psi(y)(C) = 0\} = \{y \in A^n : \beta(C, y) = 0\}.$$

The size condition on C^\perp follows from the size condition on $(\widehat{A}^n : C)$ and ψ being an isomorphism. The standard proof of the MacWilliams identities using the Poisson summation formula applies; see (Woo09) (or other standard sources) for details.

Suppose C is a left linear code and that $y \in C^\perp$, so that $\beta(C, y) = 0$. Consider ry , for $r \in R$. Then $\beta(C, ry) = \beta(\varepsilon^{-1}(r)C, y)$, by Theorem 3.4. (This particular property follows from P2.) Because C is a left R -submodule, $\varepsilon^{-1}(r)C \subset C$. Then $\beta(\varepsilon^{-1}(r)C, y) = 0$, since $y \in C^\perp$. Thus $ry \in C^\perp$, and C^\perp is a left R -submodule.

The double dual statement uses P3 in an essential way. We first show that $C \subset (C^\perp)^\perp$. To that end, suppose $x \in C$ and $y \in C^\perp$. In order that $x \in (C^\perp)^\perp$, we must show that $\beta(y, x) = 0$. But $\beta(y, x) = \beta(ex, y) = \beta(x, \varepsilon(e)y) = 0$, since $\varepsilon(e)y \in C^\perp$ because C^\perp is a left R -submodule. Note that in this derivation, Theorem 3.4 was used twice, using properties that follow from P2 and P3. Once we know that $C \subset (C^\perp)^\perp$, equality follows from the size condition (applied to C and C^\perp). \square

Remark 4.2. Because of the property $\beta(x, y) = \beta(ey, x)$, which uses P3, C^\perp is also equal to $\{x \in A^n : \beta(x, C) = 0\}$, provided the code C is linear. (The assumption of C being linear is not needed if $e = 1$.)

Remark 4.3. Although we do not include it here, the MacWilliams identities are also valid for the complete weight enumerator. See (NRS06) or (Woo09) for details.

5 Property P1—Examples

In this section we provide a number of examples of finite rings that satisfy Property P1, i.e., rings that admit an anti-isomorphism ε . Because Frobenius rings will play a prominent role in Section 6, we also mention whether the examples are Frobenius rings. (See (Lam99), (Woo99), or (Woo09) for more information about Frobenius rings.) Verifications will be left to the reader or to cited references.

Example 5.1. The first example is a non-example, i.e., an example of a finite ring that does not admit an anti-isomorphism. Let R be the ring

$$R = \begin{pmatrix} \mathbb{Z}/2^k\mathbb{Z} & \mathbb{Z}/2\mathbb{Z} \\ 0 & \mathbb{Z}/2\mathbb{Z} \end{pmatrix},$$

where $k \geq 2$. One can show that the left and right annihilators of the set of nilpotent elements are different, which would contradict the existence of an anti-isomorphism. This example is due to H. W. Lenstra, Jr., and details can be found in (Lam03, Ex. 1.22B). This ring is not Frobenius.

Example 5.2. If R is a finite commutative ring, then any ring isomorphism, in particular the identity isomorphism, is an anti-isomorphism. Some finite commutative rings (such as finite fields and Galois rings) are Frobenius; others are not.

Example 5.3. The product of rings satisfying P1 is another ring satisfying P1. That is, if finite rings R_1, \dots, R_n admit anti-isomorphisms $\varepsilon_1, \dots, \varepsilon_n$, respectively, then $R = R_1 \times \dots \times R_n$ admits the anti-isomorphism $\varepsilon = \varepsilon_1 \times \dots \times \varepsilon_n$. If every ε_i is an involution, then so is ε . The product R is Frobenius if and only if each R_i is Frobenius.

Example 5.4. Suppose a finite ring S satisfies P1 with anti-isomorphism ϵ , and suppose G is a finite group. Then the group ring $R = S[G]$ satisfies P1 with anti-isomorphism ε given by

$$\varepsilon\left(\sum_{g \in G} s_g g\right) = \sum_{g \in G} \epsilon(s_g) g^{-1}.$$

If ϵ is an involution, then so is ε . See (Lam03, Ex. 6.10) for details. If S is Frobenius, then so is $S[G]$.

Example 5.5. Suppose a finite ring S satisfies P1 with anti-isomorphism ϵ . Then the matrix ring $R = M_n(S)$ satisfies P1 with anti-isomorphism ε given by

$$\varepsilon(M) = (\epsilon(M))^T, \quad M \in M_n(S);$$

i.e., the transpose of the matrix obtained from M by taking ϵ of each entry. If ϵ is an involution, then so is ε . For details, see (Lam03, Ex. 1.22). If S is Frobenius, so is $M_n(S)$.

Example 5.6. Suppose a finite ring S satisfies P1 with anti-isomorphism ϵ . Then the upper triangular matrix ring $U_n(S)$ and the lower triangular matrix ring $L_n(S)$ satisfy P1 with anti-isomorphism ε given by:

$$\varepsilon(M)_{i,j} = \epsilon(M_{n+1-j, n+1-i}), \quad M \in U_n(S), L_n(S);$$

i.e., the (i, j) -entry of $\varepsilon(M)$ is ϵ of the $(n+1-j, n+1-i)$ -entry of M . If ϵ is an involution, then so is ε . See (Lam03, Ex. 1.22) for more details. For $n \geq 2$, the rings $U_n(S)$ and $L_n(S)$ are not Frobenius.

Example 5.7. This example involves certain finite quotients of skew polynomial rings. Let \mathbb{F}_q be a finite field of characteristic p and order $q = p^f$. Suppose σ is an automorphism of the field \mathbb{F}_q . (Then σ necessarily fixes the prime subfield $\mathbb{F}_p \subset \mathbb{F}_q$.) Define the skew-polynomial ring $\mathbb{F}_q[X; \sigma]$, as an abelian group, to be the polynomials in X over \mathbb{F}_q , but with the ring multiplication determined by the relation

$$Xa = \sigma(a)X, \quad a \in \mathbb{F}_q.$$

If σ is not the identity automorphism, then $\mathbb{F}_q[X; \sigma]$ is a non-commutative ring. It is known that one-sided division algorithms are valid in $\mathbb{F}_q[X; \sigma]$, and thus that every left (resp., right) ideal is principal. We note that (X^{l+1}) is a two-sided ideal.

Let $R = \mathbb{F}_q[X; \sigma]/(X^{l+1})$. Then R is a finite chain ring of order q^{l+1} . Every left (resp., right) ideal is two-sided, and the ideals are

$$(1) \supset (X) \supset (X^2) \supset \dots \supset (X^l) \supset (X^{l+1}) = 0.$$



Also, R is a vector space over \mathbb{F}_q with basis $1, X, X^2, \dots, X^l$. Just as for $\mathbb{F}_q[X; \sigma]$, if σ is not the identity, then R is non-commutative. Because R is a chain ring, it is a Frobenius ring.

Theorem 5.8. *The ring $R = \mathbb{F}_q[X; \sigma]/(X^{l+1})$ admits an anti-isomorphism if and only if the field automorphism σ is an involution; i.e., σ^2 is the identity. Moreover, when σ is an involution, R admits an involution.*

Proof. Observe that every element $r \in R$ was a unique representative of the form $r = \sum_{i=0}^l r_i X^i$, with $r_i \in \mathbb{F}_q$.

Suppose σ is an involution, so that σ^2 is the identity. We seek to define $\varepsilon : R \rightarrow R$ in such a way that $\varepsilon(a) = a$ for $a \in \mathbb{F}_q$ and $\varepsilon(X) = X$. Then, for any $r = \sum_{i=0}^l r_i X^i \in R$, we would need to have

$$\varepsilon(r) = \sum_{i=0}^l \varepsilon(r_i X^i) = \sum_{i=0}^l \varepsilon(X)^i \varepsilon(r_i) = \sum_{i=0}^l X^i r_i = \sum_{i=0}^l \sigma^i(r_i) X^i.$$

Consequently, we *define* $\varepsilon : R \rightarrow R$ by this formula. Because σ is a field automorphism, ε is a homomorphism of abelian groups. Because σ is an involution, one sees that ε^2 is the identity, so that ε is an isomorphism of abelian groups.

It remains to verify that $\varepsilon(rs) = \varepsilon(s)\varepsilon(r)$, for $r, s \in R$. Write $r = \sum_{i=0}^l r_i X^i$ and $s = \sum_{j=0}^l s_j X^j$. Calculations show that the coefficient of X^k in $\varepsilon(rs)$ is

$$\sigma^k\left(\sum_{i+j=k} r_i \sigma^i(s_j)\right) = \sum_{i+j=k} \sigma^k(r_i) \sigma^{k+i}(s_j),$$

while the coefficient of X^k in $\varepsilon(s)\varepsilon(r)$ is

$$\sum_{j+i=k} \sigma^j(s_j) \sigma^{j+i}(r_i) = \sum_{j+i=k} \sigma^j(s_j) \sigma^k(r_i).$$

In these formulas (which take place in \mathbb{F}_q), the factors of $\sigma^k(r_i)$ agree. As for the factors involving s_j , observe that $\sigma^{k+i}(s_j) = \sigma^{j+2i}(s_j)$, because $i + j = k$. The latter equals $\sigma^j(s_j)$, because σ^2 is the identity. Thus ε is an involution.

For the converse, suppose $\varepsilon : R \rightarrow R$ is an anti-isomorphism. Because anti-isomorphisms map left ideals to right ideals, and vice versa, they map two-sided ideals to two-sided ideals. Because ε is an isomorphism of abelian groups, it must preserve sizes of ideals. We conclude that ε maps each ideal (X^k) to itself. In particular, $\varepsilon(X) \in (X)$, so that $\varepsilon(X) = b_1 X + b_2 X^2 + \dots + b_l X^l$, for some $b_1, \dots, b_l \in \mathbb{F}_q$, with $b_1 \neq 0$. (If $b_1 = 0$, then ε would map (X) into (X^2) and sizes of ideals would not be preserved.) Similarly, the expression for $\varepsilon(X^k) = \varepsilon(X)^k$ has lowest degree terms in degree k . (The coefficient of X^k would be $b_1 \sigma(b_1) \sigma^2(b_1) \dots \sigma^{k-1}(b_1)$, which is nonzero.)

Let $\gamma \in \mathbb{F}_q$ be a primitive element of the field \mathbb{F}_q ; i.e., γ is a generator of the multiplicative group of \mathbb{F}_q . Then $\varepsilon(\gamma) = \sum_{i=0}^l c_i X^i$, for some $c_i \in \mathbb{F}_q$. Observe that $c_0 \neq 0$, lest ε map all of R into (X) . Given the expressions for $\varepsilon(\gamma)$ and $\varepsilon(X)$ above, one could write down a formula for $\varepsilon(r)$ for any $r = \sum_{i=0}^l r_i X^i \in R$. In such a formula, the constant term of $\varepsilon(r)$ involves only the constant term of $\varepsilon(r_0)$. More specifically, if $r_0 = 0$, then the constant term of $\varepsilon(r)$ is also 0. If $r_0 \neq 0$, then

$r_0 = \gamma^t$, for some t . In that case one calculates that the constant term of $\varepsilon(r)$ is c_0^t (where c_0 is the constant term of $\varepsilon(\gamma)$).

Claim 1: c_0 is a primitive element of \mathbb{F}_q . The anti-isomorphism ε is, in particular, surjective. Thus any non-zero element a of \mathbb{F}_q must appear as the constant term of $\varepsilon(r)$ for some $r \in R$. But that implies that $a = c_0^t$ for some t , which means that c_0 is a primitive element of \mathbb{F}_q .

Remember that σ is a field automorphism of \mathbb{F}_q and that $q = p^f$. Then σ has the form $\sigma(a) = a^{p^g}$, $a \in \mathbb{F}_q$, for some $g = 0, 1, \dots, f-1$. Observe that the constant term of $\varepsilon(\sigma(\gamma)) = \varepsilon(\gamma^{p^g}) = \varepsilon(\gamma)^{p^g}$ is $c_0^{p^g} = \sigma(c_0)$.

Claim 2: $\sigma^2(c_0) = c_0$. To see this, set $u = X$ and $v = \gamma X^{l-1}$. Then $uv = \sigma(\gamma)X^l$, and a computation shows that $\varepsilon(uv) = \varepsilon(X)^l \varepsilon(\sigma(\gamma)) = \varepsilon(X)^l \sigma(c_0) = \sigma^{l+1}(c_0) \varepsilon(X)^l$. (We make use of the fact that $X^{l+1} = 0$ in R to simplify expressions.) On the other hand, a similar computation shows that $\varepsilon(v)\varepsilon(u) = \sigma^{l-1}(c_0) \varepsilon(X)^l$. Since the coefficient of X^l in $\varepsilon(X)^l$, namely $b_1 \sigma(b_1) \cdots \sigma^{l-1}(b_1)$, is nonzero, we conclude that $\sigma^{l+1}(c_0) = \sigma^{l-1}(c_0)$. Because σ is an automorphism, hence invertible, it follows that $\sigma^2(c_0) = c_0$.

Claim 3: σ is an involution. This follows immediately from the previous claims. Indeed, σ^2 is a field automorphism that fixes a primitive element of the field. Thus σ^2 fixes everything; i.e., σ^2 equals the identity. \square

Example 5.9. Our last example has its origins in algebraic topology, where it appears as a finite sub Hopf algebra of the mod 2 Steenrod algebra. There are many finite sub Hopf algebras in the Steenrod algebra, and the one we present is one of the smallest and simplest. A reference is (Whi78, Chapter VIII).

Define $\mathcal{A}(1)$ to be the \mathbb{F}_2 -algebra with 1 generated by two elements, traditionally denoted Sq^1 and Sq^2 (examples of Steenrod squares), with relations $Sq^1 Sq^1 = 0$ and $Sq^2 Sq^2 = Sq^1 Sq^2 Sq^1$. It then follows that $\mathcal{A}(1)$ has dimension 8 as a vector space over \mathbb{F}_2 , with vector space basis $1, Sq^1, Sq^2, Sq^1 Sq^2, Sq^2 Sq^1, Sq^2 Sq^2 (= Sq^1 Sq^2 Sq^1), Sq^2 Sq^1 Sq^2$, and $Sq^2 Sq^2 Sq^2$. There is an involution on $\mathcal{A}(1)$, traditionally denoted χ , such that $\chi(Sq^1) = Sq^1$ and $\chi(Sq^2) = Sq^2$. Then $\chi(Sq^1 Sq^2) = Sq^2 Sq^1$, $\chi(Sq^2 Sq^1) = Sq^1 Sq^2$, and χ fixes the other basis elements. In Figure 1, the basis elements are represented by dots, multiplication by Sq^1 is represented by a straight line segment, and multiplication by Sq^2 is represented by an arc. The involution χ flips the figure top to bottom. The ring $\mathcal{A}(1)$ is a Frobenius ring.

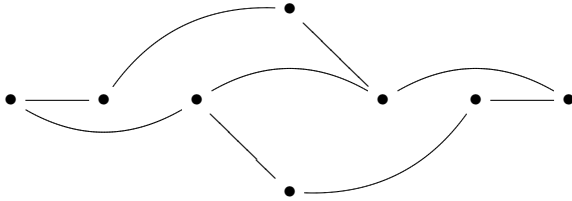


Figure 1 The ring $\mathcal{A}(1)$.

Remark 5.10. We see from these examples that Property P1 is independent of a ring being Frobenius.

6 Property P2 and Frobenius rings

In this section we suppose a finite ring R satisfies property P1, with anti-isomorphism ε . We are interested in finding examples of finite left R -modules satisfying property P2; that is, a module with an isomorphism $\psi : \varepsilon(A) \rightarrow \widehat{A}$ of right R -modules.

Theorem 6.1. *Let R be a ring with anti-isomorphism ε . Then there exists a finite R -module A satisfying property P2.*

Proof. Both the character functor $A \mapsto \widehat{A}$ and the ε -functor $A \mapsto \varepsilon(A)$ map the set of simple left R -modules bijectively to the set of simple right R -modules. Set A to be the direct sum of (one representative of each isomorphism class of) all the simple left R -modules. Then both \widehat{A} and $\varepsilon(A)$ are isomorphic to the sum of all the right simple R -modules. \square

Lemma 6.2. *Suppose R satisfies P1 with anti-isomorphism ε . Consider the left regular module ${}_R R$. Then $\varepsilon({}_R R) \cong R_R$, as right R -modules.*

Proof. Observe that ε^{-1} provides the desired isomorphism $\varepsilon({}_R R) \rightarrow R_R$. \square

Theorem 6.3. *Suppose R is a finite ring that admits an anti-isomorphism ε . Let the left module A be the ring itself: $A = {}_R R$. Then $A = {}_R R$ has property P2, $\varepsilon(A) \cong \widehat{A}$, if and only if the ring R is a finite Frobenius ring.*

Moreover, when R is a finite Frobenius ring with generating character $\varrho : R \rightarrow \mathbb{Q}/\mathbb{Z}$, an isomorphism $\psi : \varepsilon({}_R R) \rightarrow \widehat{R}_R$ is given by $\psi(b) = \varrho\varepsilon^{-1}(b)$, the right scalar multiple of $\varrho \in \widehat{R}$ by $\varepsilon^{-1}(b) \in R$. The form $\beta : R^n \times R^n \rightarrow \mathbb{Q}/\mathbb{Z}$ associated to ψ by (3.4) is given by

$$\beta(x, y) = \sum_{i=1}^n \varrho(\varepsilon^{-1}(y_i)x_i),$$

for $x = (x_1, \dots, x_n), y = (y_1, \dots, y_n) \in R^n$.

Proof. By Lemma 6.2, we know that $R_R \cong \varepsilon({}_R R)$. Thus, $A = {}_R R$ satisfies P2 if and only if $R_R \cong \varepsilon({}_R R) \cong \widehat{R}_R$. But $R_R \cong \widehat{R}_R$ if and only if R is Frobenius, by (Woo99, Theorem 3.10).

Assume R is Frobenius, so that $R_R \cong \widehat{R}_R$. This isomorphism implies the existence of a character $\varrho \in \widehat{R}_R$ (called a *generating character*) so that $r \mapsto \varrho r$ is the isomorphism $R_R \rightarrow \widehat{R}_R$. By the proof of Lemma 6.2, we know that ε^{-1} provides an isomorphism $\varepsilon({}_R R) \rightarrow R_R$. Thus $\psi(b) = \varrho\varepsilon^{-1}(b)$ is the composition of these isomorphisms $\varepsilon({}_R R) \rightarrow R_R \rightarrow \widehat{R}_R$. The formula for β now follows from (3.4). \square

Example 6.4. For all of the Frobenius rings appearing in the examples in Section 5, the left module $A = {}_R R$ satisfies P2.

Example 6.5. Let $R = \mathcal{A}(1)$ be the 8-dimensional algebra over \mathbb{F}_2 of Example 5.9; $\mathcal{A}(1)$ is a subalgebra of the mod 2 Steenrod algebra. Algebraic topology provides a rich source of $\mathcal{A}(1)$ -modules, because the mod 2 cohomology of any finite CW complex is a finite module over the Steenrod algebra and, hence, by restriction of scalars, a finite module over $\mathcal{A}(1)$. (See (Whi78, Chapter VIII).)

Here is a specific example: the mod 2 cohomology of the real projective space $\mathbb{R}P^n$. It is known that $H^*(\mathbb{R}P^n; \mathbb{F}_2)$ is a truncated polynomial algebra:

$$H^*(\mathbb{R}P^n; \mathbb{F}_2) \cong \mathbb{F}_2[a]/(a^{n+1}),$$

with $a \in H^1(\mathbb{R}P^n; \mathbb{F}_2)$, $a \neq 0$. The actions of the generators Sq^1, Sq^2 of $\mathcal{A}(1)$ are

$$Sq^i(a^j) = \begin{cases} \binom{j}{i} a^{i+j}, & i+j \leq n, \\ 0, & i+j > n, \end{cases}$$

where the binomial coefficient is calculated mod 2. See (Whi78, p. 400).



Figure 2 The $\mathcal{A}(1)$ -module $H^*(\mathbb{R}P^7; \mathbb{F}_2)$.

Write $M = H^*(\mathbb{R}P^n; \mathbb{F}_2)$; M is both a left $\mathcal{A}(1)$ -module and a vector space over \mathbb{F}_2 of dimension $n + 1$. A basis of M is $1 = a^0, a, a^2, \dots, a^n$. Figure 2 displays $H^*(\mathbb{R}P^7; \mathbb{F}_2)$ as an $\mathcal{A}(1)$ -module, with the basis elements being represented by dots and Sq^1 and Sq^2 being displayed using the same conventions as in Figure 1. The character module \widehat{M} is both a right $\mathcal{A}(1)$ -module and an \mathbb{F}_2 -vector space, with dual basis denoted $\varpi_0, \varpi_1, \dots, \varpi_n$. As characters, $\varpi_i : M \rightarrow \mathbb{Q}/\mathbb{Z}$ satisfy $\varpi_i(a^j) = (1/2)\delta_{i,j} \in \mathbb{Q}/\mathbb{Z}$, where δ is the Kronecker delta.

One calculates in the right $\mathcal{A}(1)$ -module \widehat{M} that

$$\varpi_i Sq^j = \begin{cases} \binom{i-j}{j} \varpi_{i-j}, & i-j \geq 0, \\ 0, & i-j < 0. \end{cases}$$

Define $\psi : M \rightarrow \widehat{M}$ by $\psi(a^j) = \varpi_{n-j}$. Then one calculates that

$$\psi(a^j Sq^i) = \psi(\chi(Sq^i a^j)) = \psi(Sq^i a^j) = \psi\left(\binom{j}{i} a^{i+j}\right) = \binom{j}{i} \varpi_{n-i-j},$$

while

$$\varpi_{n-j} Sq^i = \binom{n-j-i}{i} \varpi_{n-j-i}.$$

The reader will verify that the binomial coefficients agree mod 2 when $n = 4l + 3$. Thus $M = H^*(\mathbb{R}P^n; \mathbb{F}_2)$ satisfies P2 when $n = 4l + 3$.

Both M and \widehat{M} have the same displays as in Figure 2, but with Sq^1 and Sq^2 going from left to right for M , while going from right to left for \widehat{M} . Then M will satisfy P2 when its display is left-right symmetric, which happens when $n = 4l + 3$.

7 Property P3 and self-dual codes

In this final section we discuss property P3 and offer some examples of self-dual codes. The investigation of these codes is in its infancy, and all the examples are of length 1. More research will be needed in order to produce better examples.

Remember that a finite Frobenius ring R satisfies $R_R \cong \widehat{R}_R$, (Woo99, Theorem 3.10). Thus there exists a character $\varrho \in \widehat{R}$ called a *generating character* such that $R_R \rightarrow \widehat{R}_R$, $r \mapsto \varrho r$ (right scalar multiplication), is an isomorphism of right R -modules.

Lemma 7.1. *Suppose R is a Frobenius ring with generating character ϱ . If R satisfies P1 with anti-isomorphism ε , then there exists a unit $e \in R$ such that $\varrho \circ \varepsilon = \varrho e$. That is,*

$$\varrho(\varepsilon(r)) = (\varrho e)(r) = \varrho(er), \quad r \in R.$$

Proof. We make use of a result (Woo99, Lemma 4.1, Theorem 4.3) that says that a character $\varpi \in \widehat{R}$ is a generating character if and only if $\ker \varpi$ contains no nonzero left (resp., right) ideals.

Claim: $\varrho \circ \varepsilon$ is a generating character. Suppose I is a left ideal with $I \subset \ker(\varrho \circ \varepsilon)$. Then $\varepsilon(I)$ is a right ideal with $\varepsilon(I) \subset \ker \varrho$. Because ϱ is a generating character, $\varepsilon(I) = 0$. But ε is bijective, so $I = 0$, too. Thus $\ker(\varrho \circ \varepsilon)$ contains no nonzero left ideals, and we conclude that $\varrho \circ \varepsilon$ is a generating character.

Both ϱ and $\varrho \circ \varepsilon$ are generators of \widehat{R}_R , so they are scalar multiples of each other. By a result of Bass (Bas64, Lemma 6.4), they must be unit multiples of each other. Thus there exists a unit $e \in R$ such that $\varrho \circ \varepsilon = \varrho e$. \square

In the statement of the next theorem, we use Theorem 6.3 and Lemma 7.1.

Theorem 7.2. *Suppose R is a finite Frobenius ring with generating character ϱ . Suppose R satisfies P1 with anti-isomorphism ε and $A = {}_R R$ satisfies P2 with isomorphism $\psi : \varepsilon({}_R R) \rightarrow \widehat{R}_R$ given by $\psi(b) = \varrho \varepsilon^{-1}(b)$ (right scalar multiplication).*

Suppose ε is an involution and that $\varrho \circ \varepsilon = \varrho e$ with unit e being central (i.e., e is in the center of R ; it commutes with every element of R). Then ψ satisfies P3, i.e., $\widehat{\psi} = \psi e$.

Proof. Remember that $\widehat{\psi}(a)(b) = \psi(b)(a)$, $a, b \in R$. Because ε is an involution, $\psi(b) = \varrho \varepsilon(b)$. Then the result follows from the following computation.

$$\begin{aligned} \widehat{\psi}(a)(b) &= \psi(b)(a) = (\varrho \varepsilon(b))(a) = \varrho(\varepsilon(b)a) = \varrho(\varepsilon(\varepsilon(a)b)) \\ &= (\varrho \circ \varepsilon)(\varepsilon(a)b) = (\varrho e)(\varepsilon(a)b) = \varrho(e\varepsilon(a)b) = \varrho(\varepsilon(a)eb) \\ &= (\varrho \varepsilon(a))(eb) = \psi(a)(eb) = (\psi(a)e)(b). \end{aligned} \quad \square$$

We conclude this section with several examples.

Example 7.3. Let \mathbb{F}_q be a finite field of order $q = p^f$. Then \mathbb{F}_q is a Frobenius ring with generating character ϑ , as follows. Let $\text{tr} : \mathbb{F}_q \rightarrow \mathbb{F}_p$ be the trace map from \mathbb{F}_q to its prime subfield. If we view elements of \mathbb{F}_p as integers mod p , then $\vartheta(a) = \text{tr}(a)/p \in \mathbb{Q}/\mathbb{Z}$. See (Woo99, Example 4.4(i)).

Let G be a finite group (with unit element e), and let $R = \mathbb{F}_q[G]$ be the group ring of G over \mathbb{F}_q . Let $r = \sum_{g \in G} r_g g \in R$, where $r_g \in \mathbb{F}_q$. Define $\varrho \in \widehat{R}$ by $\varrho(r) = \vartheta(r_e)$, where $r_e \in \mathbb{F}_q$ is the coefficient of e in $r \in R$. Then ϱ is a generating character of R (Woo99, Example 4.4(v)). By Example 5.4 and Theorem 6.3, $\beta : R \times R \rightarrow \mathbb{Q}/\mathbb{Z}$ has the form

$$\beta(r, s) = \varrho\left(\left(\sum_{h \in G} s_h h^{-1}\right)\left(\sum_{g \in G} r_g g\right)\right) = \vartheta\left(\sum_{g \in G} s_g r_g\right),$$

for $r, s \in R$. Thus, $\beta(r, s) = \beta(s, r)$, and P3 is satisfied.

Example 7.4. Let Σ_3 be the symmetric group on three letters; $|\Sigma_3| = 6$. The elements of Σ_3 are denoted $e, \sigma, \sigma^2, \tau, \tau\sigma, \tau\sigma^2$, with $\tau^2 = e$, $\sigma^3 = e$, and $\sigma\tau = \tau\sigma^2$. Let $R = \mathbb{F}_2[\Sigma_3]$ be the group algebra of Σ_3 over \mathbb{F}_2 , and let

$$\begin{aligned} e_1 &= e + \sigma + \sigma^2, \\ e_2 &= e + \sigma + \tau\sigma + \tau\sigma^2, \\ e_3 &= e + \sigma^2 + \tau\sigma + \tau\sigma^2. \end{aligned}$$

Then e_1, e_2, e_3 are orthogonal idempotents that sum to e , which is the multiplicative identity of R .

The left ideals Re_2 and Re_3 are isomorphic, and they are simple, with dimension 2 over \mathbb{F}_2 (other basis elements are τe_2 and τe_3 , respectively). The left ideal Re_1 is indecomposable and of dimension 2 over \mathbb{F}_2 , but it is not simple. It has a 1-dimensional subideal $R(e + \tau)e_1$. The left ideals $C_1 = R(e + \tau)e_1 + Re_2$ and $C_2 = R(e + \tau)e_1 + Re_3$ are examples of self-dual codes in R .

Example 7.5. Let $R = \mathbb{F}_q[X; \sigma]/(X^{l+1})$. Let ϑ be a generating character for \mathbb{F}_q , as in Example 7.3. Let $r = \sum_{i=0}^l r_i X^i \in R$, where $r_i \in \mathbb{F}_q$. Define $\varrho : R \rightarrow \mathbb{Q}/\mathbb{Z}$ by $\varrho(r) = \vartheta(r_l)$, where r_l is the coefficient of X^l in $r \in R$. Then ϱ is a generating character of R . Here is the argument. The ring R is a chain ring, so (X^l) is a minimal ideal (and (X^l) is the socle of R). If $\ker \varrho$ were to contain a nonzero ideal, then $(X^l) \subset \ker \varrho$. That is, $\varrho((X^l)) = 0$. But $\varrho(r_l X^l) = \vartheta(r_l)$, which is not identically zero, because ϑ is a generating character of \mathbb{F}_q .

Now assume that σ is an involution, so that R admits an involution ε and $A = R$ satisfies P2 (by Theorems 5.8 and 6.3). By Example 5.7, $\beta : R \times R \rightarrow \mathbb{Q}/\mathbb{Z}$ has the form

$$\begin{aligned} \beta(r, s) &= \varrho(\varepsilon(s)r) = \varrho\left(\left(\sum_{i=0}^l \sigma^i(s_i)X^i\right)\left(\sum_{j=0}^l r_j X^j\right)\right) \\ &= \varrho\left(\sum_{i,j} \sigma^i(s_i)\sigma^i(r_j)X^{i+j}\right) = \vartheta\left(\sum_{i+j=l} \sigma^i(s_i)\sigma^i(r_j)\right) \\ &= \sum_{i+j=l} \vartheta(\sigma^i(s_i r_j)) = \sum_{i+j=l} \vartheta(s_i r_j). \end{aligned}$$

Here, we have used the fact that $\vartheta(\sigma(a)) = \vartheta(a)$ for $a \in \mathbb{F}_q$, because the trace satisfies $\text{tr}(\sigma(a)) = \text{tr}(a)$, $a \in \mathbb{F}_q$. Because the formula above for β is symmetric in r, s , we see that $\beta(r, s) = \beta(s, r)$, and P3 is satisfied.

When $l + 1 = 2k$ is even, $C = (X^k)$ is a self-dual code.

Example 7.6. Let $R = \mathcal{A}(1)$ be the 8-dimensional \mathbb{F}_2 algebra of Example 5.9. Refer to the \mathbb{F}_2 -vector space basis elements as follows: $b_0 = 1, b_1 = \text{Sq}^1, b_2 = \text{Sq}^2, b_3 = \text{Sq}^1 \text{Sq}^2, b_{3'} = \text{Sq}^2 \text{Sq}^1, b_4 = \text{Sq}^2 \text{Sq}^2, b_5 = \text{Sq}^2 \text{Sq}^1 \text{Sq}^2$, and $b_6 = \text{Sq}^2 \text{Sq}^2 \text{Sq}^2$. Set $I = \{0, 1, 2, 3, 3', 4, 5, 6\}$. We will write a typical element of $\mathcal{A}(1)$ as $r = \sum_{i \in I} r_i b_i$. Let ϑ be a generating character for \mathbb{F}_2 ; if we view elements of \mathbb{F}_2 as integers mod 2, then $\vartheta(a) = a/2 \in \mathbb{Q}/\mathbb{Z}$. Define $\varrho \in R$ by $\varrho(r) = \vartheta(r_6)$, where r_6 for the coefficient of b_6 in $r \in \mathcal{A}(1)$. Then ϱ is a generating character of $\mathcal{A}(1)$. Indeed, the socle

of $\mathcal{A}(1)$ is the simple 1-dimensional ideal generated by $b_6 = \text{Sq}^2 \text{Sq}^2 \text{Sq}^2$, and the argument given in Example 7.5 applies.

By Example 5.9, $\beta : \mathcal{A}(1) \times \mathcal{A}(1) \rightarrow \mathbb{Q}/\mathbb{Z}$ has the form

$$\beta(r, s) = \varrho(\chi(s)r) = \vartheta\left(\sum_{i+j=6} s_i r_j\right),$$

were we agree that $3 + 3' = 6$ (and $3 + 3$ and $3' + 3'$ do not sum to 6). Since the sum above is symmetric in r, s , we have $\beta(r, s) = \beta(s, r)$, and P3 is satisfied.

The left ideal $C = \mathcal{A}(1)(\text{Sq}^1 \text{Sq}^2 + \text{Sq}^2 \text{Sq}^1)$ is a self-dual code. Figures 3 and 4 offer two displays of this code. The filled-in dots are a basis for the code; the open dots are to allow comparison with Figure 1. In Figure 3, the symbol \circlearrowleft is used because the basis elements $\text{Sq}^1 \text{Sq}^2$ and $\text{Sq}^2 \text{Sq}^1$ are added.

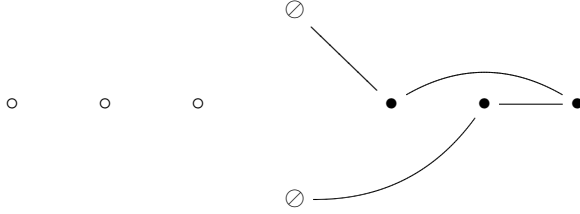


Figure 3 An $\mathcal{A}(1)$ -linear self-dual code in $\mathcal{A}(1)$: first view.

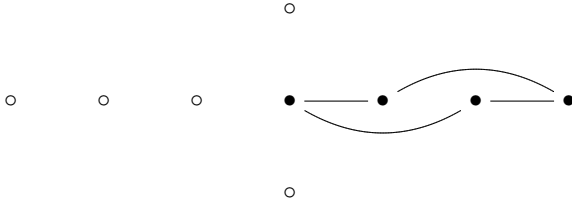


Figure 4 An $\mathcal{A}(1)$ -linear self-dual code in $\mathcal{A}(1)$: second view.

Example 7.7. Let $R = \mathcal{A}(1)$, as in Example 7.6. Let $M = H^*(\mathbb{R}P^n; \mathbb{F}_2)$, with $n = 4l + 3$, as in Example 6.5; M satisfies P2. A typical element of M has the form $r = \sum_{i=0}^n r_i a^i$. Using ψ defined in Example 6.5, $\beta : M \times M \rightarrow \mathbb{Q}/\mathbb{Z}$ has the form

$$\beta(r, s) = \psi(s)(r) = \left(\sum_{i=0}^n s_i \varpi_{n-i}\right) \left(\sum_{j=0}^n r_j a^j\right) = \sum_{i,j} s_i r_j \varpi_{n-i}(a^j) = \vartheta\left(\sum_{i=0}^n s_i r_{n-i}\right),$$

where ϑ is the generating character for \mathbb{F}_2 . The formula is symmetric in r, s , so $\beta(r, s) = \beta(s, r)$, and P3 is satisfied.

The vector subspace C spanned by $a^{2l+2}, a^{2l+3}, \dots, a^{4l+3}$ has dimension $2l + 2$ and is a left $\mathcal{A}(1)$ -submodule of M ; C is a self-dual code. Figure 5 displays this self-dual code when $n = 7$. The filled-in dots are a basis for the code; the open dots are to allow comparison with Figure 2.





Figure 5 An $\mathcal{A}(1)$ -linear self-dual code in $H^*(\mathbb{R}P^7; \mathbb{F}_2)$.

Remark 7.8. The examples in this section have the feature that the rings R are algebras over a finite field \mathbb{F}_q , and the R -modules are vector spaces over \mathbb{F}_q . Then R -linear self-dual codes can be viewed as self-dual codes over \mathbb{F}_q with additional symmetry coming from R . One caution: the self-duality over \mathbb{F}_q may involve an inner product β different from the standard dot product. The standard dot product occurs for group rings. An alternating form occurs in the other examples.

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