Isometry Groups of Additive Codes over Finite Fields

Jay A. Wood

Department of Mathematics
Western Michigan University
http://homepages.wmich.edu/∼jwood

Central China Normal University
Wuhan, Hubei
May 15, 2018

Isometries of additive codes over finite fields

- Additive codes as linear codes over modules
- Failure of EP
- Monomial and isometry groups
- Examples
- Criteria in terms of multiplicity functions
- Structure of ker W
- EP for short codes
- Building codes with prescribed groups
- Extreme examples

Additive $\mathbb{F}_4$-codes

- There has been interest in additive codes with alphabet $A = \mathbb{F}_4$.
- Such codes are the same as R-linear codes over A with $R = \mathbb{F}_2$ and $A = \mathbb{F}_4$, regarding $\mathbb{F}_4$ as an $\mathbb{F}_2$-vector space of dimension 2.
- Generalize to case of $R = M_{k\times k}(\mathbb{F}_q)$ and $A = M_{k\times\ell}(\mathbb{F}_q)$. Information module will be $M = M_{k\times m}(\mathbb{F}_q)$.
- Call this the **matrix module context**.

Failure of EP when $k < \ell$

- Recall that EP for Hamming weight fails in the matrix module context when $k < \ell$ and $k < m$.
- In terms of the W-map:

$$W : F_0(\mathcal{O}^\sharp, \mathbb{Q}) \to F_0(\mathcal{O}, \mathbb{Q})$$

fails to be injective for all information modules M.

Isometry group

- General set-up: ring R, alphabet A, weight w on A.
- Let $C \subseteq A^n$ be an R-linear code.
- Consider self-maps: linear isometries $f : C \to C$; i.e., $w(cf) = w(c)$, for all $c \in C$.
- When C is given as the image of a parametrized code $\Lambda : M \to A^n$, we define the **isometry group**:

$$\mathrm{Isom}(C) = \{g \in \mathrm{GL}_R(M) : \text{there exists a linear isometry } f : C \to C \text{ such that } g\Lambda = \Lambda f\}.$$

- View isometries on M rather than C.

Monomial group

- Recall that the weight w on A has a right symmetry group $G_{\mathrm{rt}} = \{\phi \in \mathrm{GL}_R(A) : w(a\phi) = w(a), a \in A\}$.
- For linear code $C \subseteq A^n$, define the **monomial group**

$$\mathrm{Monom}(C) = \{T : A^n \to A^n, G_{\mathrm{rt}}\text{-monomial transformation, with } CT = C\}.$$

Restriction map

- Any $T \in \text{Monom}(C)$, when restricted to C, gives an isometry on C. By viewing the isometry on M, we get a group homomorphism

$$\text{restr} : \text{Monom}(C) \to \text{Isom}(C).$$

- Denote $\ker \text{restr} = \text{Monom}_0(C)$. Think of repeated columns in a generator matrix: can permute columns without changing the codewords.
- If EP holds, then restr is surjective.

Main question

- When EP fails, restr may not be surjective for all linear codes C or information modules M.
- Then $\mathrm{restr}(\mathrm{Monom}(C)) \subseteq \mathrm{Isom}(C) \subseteq \mathrm{GL}_R(M)$.
- What subgroups of $\mathrm{GL}_R(M)$ can occur as $\mathrm{restr}(\mathrm{Monom}(C))$ and $\mathrm{Isom}(C)$?

Example 1 (a)

- Additive code over $\mathbb{F}_4 = \mathbb{F}_2[\omega]/(\omega^2 + \omega + 1)$ with generator matrix G_1 and list of codewords. $M = \mathbb{F}_2^3$.

$$G_1 = \begin{bmatrix} 1 & \omega & 0 \\ \omega & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}, \qquad \begin{matrix} 0 & 0 & 0 \\ 1 & \omega & 0 \\ \omega & 1 & 0 \\ \omega^2 & \omega^2 & 0 \\ 1 & 0 & 1 \\ 0 & \omega & 1 \\ \omega^2 & 1 & 1 \\ \omega & \omega^2 & 1 \end{matrix}$$

Example 1 (b)

- Consider three elements of $\mathrm{GL}_R(M) = \mathrm{GL}(3, \mathbb{F}_2)$:

$$f_1 = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \quad f_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}, \quad f_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix}.$$

- f_1, f_2 generate restr(Monom(C)), a Klein 4-group. But f_1, f_3 generate Isom(C), a dihedral group of order 8. ($f_2 = f_1 f_3^2$.)
- Magma found only the cyclic 2-group generated by $f_1 f_2$. (Magma looks only for $\mathbb{F}_4$-linear maps.)

Example 2 (a)

- Additive code over $\mathbb{F}_4$ with generator matrix G_2 and list of codewords. Again, $M = \mathbb{F}_2^3$.

$$G_2 = \begin{bmatrix} 0 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & \omega & \omega \\ \omega & \omega & 1 & 0 & \omega^2 \end{bmatrix},$$

$$\begin{matrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & \omega & \omega \\ 1 & 1 & 0 & \omega^2 & \omega^2 \\ \omega & \omega & 1 & 0 & \omega^2 \\ \omega & \omega^2 & 0 & 1 & \omega \\ \omega^2 & \omega & 0 & \omega & 1 \\ \omega^2 & \omega^2 & 1 & \omega^2 & 0 \end{matrix}$$

Example 2 (b)

- Consider three elements of $\mathrm{GL}_R(M) = \mathrm{GL}(3, \mathbb{F}_2)$:

$$f_4 = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix}, \quad f_5 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix}, \quad f_6 = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix}.$$

- These elements generate $\mathrm{restr}(\mathrm{Monom}(C)) \cong \Sigma_4$, the symmetric group on 4 elements, while $\mathrm{Isom}(C) = \mathrm{GL}(3, \mathbb{F}_2)$, the simple group of order 168.
- Magma found only a cyclic 4-group generated by $f = f_4 f_5 f_6 f_4 f_5 f_4 f_6$.

Closure for group actions

- Some of the hypotheses of the main result involve a notion of **closure** with respect to a group action.
- This idea goes back at least to Wielandt, 1964.
- Suppose a finite group G acts on a set X.
- A subgroup $H \subseteq G$ partitions X into H-orbits.
- Define the **closure** of H with respect to the action:

$$\bar{H} = \{g \in G : g \operatorname{orb}_H(x) = \operatorname{orb}_H(x), x \in X\}.$$

- Subgroup $H \subseteq G$ is **closed** with respect to the action if $\bar{H} = H$.

Closure conditions

- Usual set-up: ring R, alphabet A, weight w, information module M. Orbit spaces $\mathcal{O}$ and $\mathcal{O}^\sharp$.
- $\mathcal{O} = G_{\mathrm{lt}}\backslash M$: $\mathrm{GL}_R(M)$ acts on the right of $\mathcal{O}$, and on the left of $F_0(\mathcal{O}, \mathbb{Q})$.
- $\mathcal{O}^\sharp = \mathrm{Hom}_R(M, A)/G_{\mathrm{rt}}$: $\mathrm{GL}_R(M)$ acts on the left, and on the right of $F_0(\mathcal{O}^\sharp, \mathbb{Q})$: $(\eta f)([\lambda]) = \eta([f\lambda])$.
- For $H_1 \subseteq H_2 \subseteq \mathrm{GL}_R(M)$, will want H_1 to be closed for the $\mathcal{O}^\sharp$-action and H_2 closed for the $\mathcal{O}$-action.
- "Not every subgroup gets to be an isometry group."

Statement of main result

Theorem

Matrix module context with $k < \ell < m$ *and the Hamming weight on* A*. For any choice of subgroups* $H_1 \subseteq H_2 \subseteq \mathrm{GL}_R(M)$ *with* H_1 *closed for the* $\mathcal{O}^\sharp$*-action and* H_2 *closed for the* $\mathcal{O}$*-action, there exists a linear code* C *modeled on* M *such that* $H_1 = \mathrm{restr}(\mathrm{Monom}(C))$ *and* $H_2 = \mathrm{Isom}(C)$*.*

Corollary

Same matrix module context. There exists a linear code C *modeled on* M *with* $\mathrm{restr}(\mathrm{Monom}(C)) = \{\mathbb{F}_q^\times \cdot \mathrm{id}_M\}$ *and* $\mathrm{Isom}(C) = \mathrm{GL}_R(M)$*.*

Using multiplicity functions

- Up to G_{rt}-monomial transformations, a parametrized code $\Lambda : M \to A^n$ is determined by its multiplicity function $\eta_\Lambda \in F_0(\mathcal{O}^\sharp, \mathbb{N})$.
- Recall the W-map: $W : F_0(\mathcal{O}^\sharp, \mathbb{Q}) \to F_0(\mathcal{O}, \mathbb{Q})$.
- Recall the right action of $\mathrm{GL}_R(M)$ on $F_0(\mathcal{O}^\sharp, \mathbb{Q})$: $(\eta f)([\lambda]) = \eta([f\lambda])$.
- For $f \in \mathrm{GL}_R(M)$, $f \in \mathrm{restr}(\mathrm{Monom}(\eta))$ if and only if $\eta f = \eta$.
- For $f \in \mathrm{GL}_R(M)$, $f \in \mathrm{Isom}(\eta)$ if and only if $\eta f - \eta \in \ker W$.

Structure of ker W (a)

- In the matrix module context, $\mathcal{O}^\sharp$ is the set of CRE matrices of size $m \times \ell$, while $\mathcal{O}$ is the set of RRE matrices of size $k \times m$.
- Remember $k < \ell < m$. By dimension counting,

$$\ker W \geq \sum_{i=k+1}^{\ell} \begin{bmatrix} m \\ i \end{bmatrix}_q , \qquad (1)$$

using the q-binomial coefficients.

Structure of ker W (b)

- The orbit space $\mathcal{O}^\sharp$ is partitioned by rank.
- By explicit constructions, one produces independent elements $\eta_{[\lambda]} \in \ker W$. For each $i = k+1, \ldots, \ell$, one produces $\left[{m \atop i}\right]_q$ of them, each $\eta_{[\lambda]}$ supported on $[\lambda]$ of rank i and on specific elements of smaller rank. ("Triangular.") This produces as many independent elements of ker W as the sum in (1).
- Separately, one shows that W is surjective, so there is equality in (1), and we have an explicit basis for ker W.

Aside: EP for short codes

- Serhii Dyshko (Toulon) has shown that EP holds even when $k < \ell$, **provided** n is sufficiently small ($n \leq q$ when $k = 1$).
- Elements of ker W affect the length of the code.
- The exact details of this need to be better understood.

Idea of proof (a)

- Elements $[x] \in \mathcal{O}$ have a well-defined rank, $\mathrm{rk}[x]$. The $\mathrm{GL}_R(M)$-action preserves this rank.
- Pick a function w on $\mathcal{O}$ that (1) is constant on each and separates the H_2-orbits on $\mathcal{O}$ and (2) is an increasing function of $\mathrm{rk}[x]$.
- Because W is surjective, there exists η with $W(\eta) = w$. A priori, η has rational values.
- Can modify η to have non-negative integer values and still satisfy (1) and (2).

Idea of proof (b)

- Replace η by an averaged version so that η is also constant on the H_2-orbits on $\mathcal{O}^\sharp$. This does not change $W(\eta)$. Clear denominators of η, which scales everything.
- At this point, η has non-negative integer values, is constant on H_2-orbits on $\mathcal{O}^\sharp$, and is constant on and separates H_2-orbits on $\mathcal{O}$.

Idea of proof (c)

- Claim $\mathrm{restr}(\mathrm{Monom}(\eta)) = \mathrm{Isom}(\eta) = H_2$.
- From η constant on H_2-orbits on $\mathcal{O}^\sharp$, $H_2 \subseteq \mathrm{restr}(\mathrm{Monom}(\eta))$.
- We always have $\mathrm{restr}(\mathrm{Monom}(\eta)) \subseteq \mathrm{Isom}(\eta)$.
- Suppose $f \in \mathrm{Isom}(\eta)$. Because $w = W(\eta)$ separates H_2-orbits on $\mathcal{O}$, $w(xf) = w(x)$ implies $f \in \bar{H}_2$. The closure hypothesis implies $f \in H_2$.

Idea of proof (d)

- Modify η using $\eta_{[\lambda]} \in \ker W$ to separate H_1-orbits on $\mathcal{O}^\sharp$ (rank-by-rank, from rank ℓ down to rank $k+1$).
- Because of "triangular" form of $\eta_{[\lambda]}$, a change at rank i does not disturb changes at higher ranks.
- The final η preserves H_1-orbits on $\mathcal{O}^\sharp$, so $H_1 \subseteq \mathrm{restr}(\mathrm{Monom}(\eta))$. Conversely, any $f \in \mathrm{restr}(\mathrm{Monom}(\eta))$ preserves H_1-orbits on $\mathcal{O}^\sharp$ (η separates), so $f \in \bar{H}_1$. Closure implies $f \in H_1$.
- Because modifications were made by $\eta_{[\lambda]} \in \ker W$, $W(\eta)$ has not changed. We still have $\mathrm{Isom}(\eta) = H_2$.

Extreme example (a)

- $R = \mathbb{F}_2$, $A = \mathbb{F}_4$, $M = \mathbb{F}_2^3$. Multiplicities as indicated. Length $n = 28$.

multiplicity	1	4	2	2	4	1	3	5	6
	1	0	0	1	1	1	1	1	1
G	0	1	1	ω	ω	ω	ω	0	1
	1	0	1	0	ω	1	ω^2	ω	ω

- All codewords have weight 22, so $\mathrm{Isom}(C) = \mathrm{GL}(3, \mathbb{F}_2)$, while $\mathrm{restr}(\mathrm{Monom}(C)) = \{\mathrm{id}_M\}$.

Extreme example (b)

- Additive code over $\mathbb{F}_9 = \mathbb{F}_3[\omega]/(\omega^2 - \omega - 1)$.

mult.	5	3	6	1	1	1	2	2	2	4	3	2
	0	0	0	1	1	1	1	1	1	1	1	1
G_3	0	1	1	0	0	0	1	1	1	-1	-1	-1
	1	1	-1	0	1	-1	0	1	-1	0	1	-1

6	3	7	8	9	6	4	5	2	3	1
1	1	1	1	1	1	1	1	1	1	1
0	1	ω	ω	ω	ω	ω	ω	ω	ω	ω
ω	ω	0	1	-1	ω	$\omega+1$	$\omega-1$	$-\omega$	$-\omega+1$	$-\omega-1$

Extreme example (b) continued

- Code has length $n = 86$; all codewords have weight 72.
- $\mathrm{Isom}(C) = \mathrm{GL}(3, \mathbb{F}_3)$, of order $11,232$.
- $\mathrm{restr}(\mathrm{Monom}(C)) = \{\pm \mathrm{id}_M\}$ is minimum possible.

Other alphabets

- Most of the result carries over to any alphabet with non-cyclic socle, such as non-Frobenius rings.
- Get $\text{restr}(\text{Monom}(\eta)) \subseteq H_1$ only, but still have $H_2 = \text{Isom}(\eta)$.
- This is enough to get the extreme cases.