

Ring Involutions and Self-Dual Codes

Jay A. Wood

Department of Mathematics
Western Michigan University
jay.wood@wmich.edu

AMS meeting, Lexington, Kentucky
March 27, 2010

Acknowledgments

- ▶ Much of the material in this talk is my interpretation of work previously done by Nebe, Rains, and Sloane, *Self-Dual Codes and Invariant Theory*, Springer, 2006.
- ▶ Results to appear in special issue dedicated to Vera Pless of Int. J. of Information and Coding Theory.

Wanted: a good notion of self-dual codes over non-commutative rings

- ▶ In general, the dual code to a left linear code will be a right linear code.
- ▶ In general, left submodules are not right submodules.
- ▶ How can we get a linear code to be self-dual?
- ▶ Will “standard properties” hold?

The classical case—finite fields

- ▶ On \mathbb{F}_q^n , the standard \mathbb{F}_q -valued dot product is a nondegenerate, symmetric bilinear form.
- ▶ If $C \subset \mathbb{F}_q^n$ is a linear code of dimension k , then the *dual code* is

$$C^\perp = \{y \in \mathbb{F}_q^n : x \cdot y = 0, \text{ all } x \in C\}.$$

- ▶ We work with Hamming weights throughout.

“Standard properties” in the classical case

- ▶ $C^\perp \subset \mathbb{F}_q^n$.
- ▶ C^\perp is a linear code.
- ▶ $\dim C + \dim C^\perp = n$; or $|C||C^\perp| = |\mathbb{F}_q^n|$.
- ▶ $(C^\perp)^\perp = C$.
- ▶ The MacWilliams identities hold.

Some notation

- ▶ Finite ring R , usually non-commutative.
- ▶ Alphabet A , a finite left R -module.
- ▶ A left *linear code* over A of length n is a left R -submodule $C \subset A^n$.

MacWilliams identities—how are they proved?

- ▶ Gleason's proof uses characters and the Poisson summation formula.
- ▶ If $C \subset A^n$ is a left linear code, then its “dual” will be the character-theoretic annihilator $(\widehat{A}^n : C)$, a right submodule of the right character module \widehat{A}^n .
- ▶ $(\widehat{A}^n : C) = \{\pi \in \widehat{A}^n : \pi(C) = 1\}$.
- ▶ Make identifications (e.g., when $A = R$, Frobenius).

Obstacles

- ▶ We want the dual code to be a submodule of A^n , not \widehat{A}^n .
- ▶ In order to have self-dual codes, we need to have the dual code be a left submodule, not a right submodule.

Making identifications, I

- ▶ Left-right identifications: assume the ring R admits an anti-isomorphism.
 - ▶ *Anti-isomorphism* $\varepsilon : R \rightarrow R$, additive isomorphism, with $\varepsilon(rs) = \varepsilon(s)\varepsilon(r)$, for all $r, s \in R$.
 - ▶ *Involution*: if $\varepsilon^2 = 1$.
- ▶ Given a left R -module M , define a right R -module $\varepsilon(M)$ to be the same additive group, but with right scalar multiplication $mr := \varepsilon(r)m$, for $m \in M$, $r \in R$.

Making identifications, II

- ▶ Characters: assume the alphabet A admits an isomorphism $\psi : \varepsilon(A) \rightarrow \widehat{A}$ of right R -modules.
- ▶ Dual code: for a left linear code $C \subset A^n$, define the *dual code* $C^\perp = \varepsilon^{-1}\psi^{-1}(\widehat{A}^n : C)$.
- ▶ The dual code is now a left submodule of A^n .

Which standard properties hold?

- ▶ $C^\perp \subset A^n$.
- ▶ C^\perp is a left linear code, if C is.
- ▶ $|C||C^\perp| = |A^n|$, from general character results.
- ▶ Double dual?—not clear. Need an extra assumption on ψ .
- ▶ MacWilliams identities hold.

Recasting in terms of biadditive form

- ▶ View characters as having values in \mathbb{Q}/\mathbb{Z} (“additive form”).
- ▶ Define $\beta : A^n \times A^n \rightarrow \mathbb{Q}/\mathbb{Z}$ by

$$\beta(x, y) = \sum \psi(y_i)(x_i), \quad x, y \in A^n.$$

Properties of β

- ▶ β is biadditive.
- ▶ β is non-degenerate.
- ▶ $\beta(rx, y) = \beta(x, \varepsilon(r)y)$, all $r \in R$, $x, y \in A^n$.
- ▶ (Extra) There exists unit $e \in R$ such that $\beta(x, y) = \beta(ey, x)$, all $x, y \in A^n$.
- ▶ $C^\perp = \{y \in A^n : \beta(C, y) = 0\}$.

An example

- ▶ G finite group. $R = F[G]$, the group algebra.
- ▶ Involution $\varepsilon(\sum a_g g) = \sum a_g g^{-1}$.
- ▶ $G = \Sigma_3$, symmetric group: $\sigma^3 = e$, $\tau^2 = e$,
 $\sigma\tau = \tau\sigma^2$.
- ▶ $C = R(e + \tau)(e + \sigma + \sigma^2) + R(e + \sigma + \tau\sigma + \tau\sigma^2)$
is a self-dual code of length 1 in $A = R$.

Questions

- ▶ Which finite rings admit anti-isomorphisms?
- ▶ Which finite modules admit isomorphisms $\psi : \varepsilon(A) \rightarrow \widehat{A}$?
- ▶ Are there interesting examples?
- ▶ Some partial results in paper in Pless issue.
- ▶ Some examples from algebraic topology: cohomology as modules over the Steenrod algebra.